



MIRROR PERESOLUTION TEST CHA NATIONAL GENERAL DE STANJARI, LIGHT

E (When Data Entered) AD-A195 508 ENTATION PAGE READ INSTRUCTIONS BEFORE COMPLETEING FORM WITH TILL TUPT 12. GOVT ACCESSION NO. 3. RECIPIENT'S CATALOG NUMBER 4. TITLE (and Subtitle) 5. TYPE OF REPORT & PERIOD COVERED 4 June 1987 to 4 June 1988 Ada Compiler Validation Summary Report: Alsys Inc., AlsyCOMP_003, V3.1, Wang PC 280 6. PERFORMING ORG. REPORT NUMBER 7. AUTHOR(s) 8. CONTRACT OR GRANT NUMBER(s) The National Computing Centre Limited Manchester, UK 9. PERFORMING ORGANIZATION AND ADDRESS 10. PROGRAM ELEMENT, PROJECT, TASK AREA & WORK UNIT NUMBERS The National Computing Centre Limited Manchester, UK 1. CONTROLLING OFFICE NAME AND ADDRESS Ada Joint Program Office 12. REPORT DATE 4 June 1987 United States Department of Defense Washington, DC 20301-3081 13. NUMBER OF PAGES 53 p. 14. MONITORING AGENCY NAME & ADDRESS(If different from Controlling Office) 15. SECURITY CLASS (of this report) The National Computing Centre Limited UNCLASSIFIED Manchester, UK 15a. DECLASSIFICATION/DOWNGRADING

16. DISTRIBUTION STATEMENT (of this Report)

Approved for public release; distribution unlimited.

17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20. If different from Report)

UNCLASSIFIED

18. SUPPLEMENTARY NOTES

N/A

19. KEYWORDS (Continue on reverse side if necessary and identify by block number)

Ada Programming language, Ada Compiler Validation Summary Report, Ada Compiler Validation Capability, ACVC, Validation Testing, Ada Validation Office, AVO, Ada Validation Facility, AVF, ANSI/MIL-STD-1815A, Ada Joint Program Office, AJPO

20. ABSTRACT (Continue on reverse side if necessary and identify by block number)

AlsyCOMP 003, V3.1. Alsys Inc., National Computing Centre Limited, Wang PC 280 under MS-DOS 3.20 (host and target). ACVC 1.8.

1473

EDITION OF 1 NOV 65 IS OBSOLETE

1 JAN 73

S/N 0102-LF-014-6601

UNCLASSIFIED

SECURITY CLASSIFICATION OF THIS PAGE (When Data Entered)

* Ta* COMPTLER
VAUTOAGION SUMMARY REPORT
Alsys Inc
Alsycomp 003. V3 1
Wang "C 280

Completion of On-site Festing
4 June 1987

Prepared By
The National Computing Centre Limited
Oxford Road
Manchester
M1 7ED
UK

Prepared For
Ada Joint Program Office
United States Department of Defense
Washington, D.C.
USA

^{*}Ada is a registered trademark of the United States Government (Ada Joint Program Office).

Ada* Compiler Validation Summary Report:

Compiler Name: AlsyCOMP 003, V3.1

Host:

Target:

Wang PC 280

under MS-DOS 3.20

Same as host

Testing Completed 4 June 1987 using ACVC 1.8

This report has been reviewed and is approved.

The National Computing Centre Ltd

Vony Gwillim Oxford Road Manchester M1 7ED

Ada Validation Office

Dr. John F. Kramer
Institute for Defense

Institute for Defense Analyses

Alexandria VA

Virginia L. Costa Ada Joint Program Office

Virginia L. Castor

Director

Department of Defense

Washington DC

Accession For

NTIS GRA&I
DTIC TAB
Unannounced
Justification

By
Distribution/
Availability Ccdes

Avail and/or
Dist Special



^{*}Ada is a registered trademark of the United States Government (Ada Joint Program Office).

EXECUTIVE SUMMARY

This Validation Summary Report (VSR) summarizes the results and conclusions of validation testing performed on the AlsyCOMP_003, V3.1 using Version 1.8 of the Ada* Compiler Validation Capability (ACVC). The AlsyCOMP_003 is hosted on a Wang PC 280 operating under MS-DOS 3.20.

On-site testing was be formed to June 1987 through 4 June 1907 at Alsys . to Maltham MA 0.154 under the direction of the National Computing [NF], according to Ada Validation Organization (AVO) policies educes. The AVF identified 2210 of the 2399 test in ACVC was one 1.8 to be processed during on-site testing of the compiler. TR Tests withdrawn at the time of validation testing, as well as 7% executable tests that make use of floating-point precision ... That supported by the implementation were not processed. Attaches 210 tests were processed, results for Class A, C, D, or E teses were examined for a react execution. Compilation listings for Class B tests were analyzed for correct diagnosis of syntax and semantic errors. Compilation and link results of Class L tests were analyzed for correct detaction of errors. There were 26 of the processed tests determined to be inapplicable; The remaining testa were passed.

The results of validation are summarized in the following table:

RESULT	CHAPTER									TOTAL			
	2	3	4	5	б	7	8	9	10	11	12	14	
Passed	102	253	334	243	161	97	136	261	129	32	218	218	2184
©.iled	0	0	0	0	0	0	0	0	0	0	0	0	0
fuapplicable	14	72	86	4	0	0	3	1	1	0	0	15	196
Withdrawn	0	5	5	0	0	1	1	2	4	O	1	0	19
ተርሞልፒ	116	330	425	247	161	98	140	264	134	32	219	233	2399

AVF concludes that these results demonstrate acceptable conformity to ANSI/MIL-STD-1815A Ada.

^{*}Ada is a registered trademark of the United States Government (Ada Joint Program Office).

		TABLE OF CONTENTS
CHAPTER	1	INTRODUCTION
ŕ	1.1 1.2 1.3 1.4	PURPOSE OF THIS VALIDATION SUMMARY REPORT
to the state of	5	CONFERRATION ENFORMATION
	2.1	CONFIGURATION TESTED
SHAPPER	3	TEST INFORMATION
	3.1 3.2 3.3 3.4 3.5 3.6 3.7 3.7.1 3.7.2 3.7.3	TEST RESULTS SUMMARY OF TEST RESULTS BY CLASS SUMMARY OF TEST RESULTS BY CHAPTER WITHDRAWN TESTS INAPPLICABLE TESTS SPLIT TESTS ADDITIONAL TESTING INFORMATION Provalidation Test Method Test Site 3-2 3-4 3-4 3-4 3-4 3-4 3-4 3-4 3-4 3-4 3-4
APPENDI	X A	COMPLIANCE STATEMENT
APPENDA	€3	APPENDIX F OF THE Ada STANDARD
APPENDI	хс	TEST PARAMETERS
APPENDI	X D	WITHDRAWN TESTS

CHAPTER L

INTRODUCTION

Main Ada compiler conforms to the Ada Standard. This report explaint all technical terms used within it and thoroughly reports the results of testing this compiler using the Ada Compiler Validation Capability (ACVC). An Ada compiler must be implemented according to the Ada Standard and any implementation-dependent features must conform to the requirements of the Ada Standard. The Ada Standard must be implemented in its entirety, and nothing can be implemented that is not in the Standard.

-dve = though all validated Ada compilers conform to the Ada Standard, it must be understood that some differences do exist implementations. The Ada Standard permits some implementation dependencies--for the maximum length of identifiers or axample, maximum values of integer types. Other differences between compilers result from characteristics of particular operating systems, hardware, implementation strategies. All of the dependencies demonstrated during the process of testing this compiler are given in this report.

PROCEEDS SERVICED REPORTED TO THE SERVICE OF SERVICED DESCRIPTION OF SERVICED

The information in this report is derived from the test results produced during validation testing. The validation process includes submitting a suite of standardized tests, the ACVC, as inputs to an Ada compiler and evaluating the results. The purpose of validating is to ensure conformity of the compiler to the Ada Standard by testing that the compiler properly implements legal language constructs and that it identifies and rejects illegal language constructs. The testing also identifies behaviour that is implementation dependent but permitted by the Ada Standard. Six classes of tests are used. These tests are designed to perform checks at compile time, at link time, and during execution.

1.1 PURPOSE OF THIS VALIDATION SUMMARY REPORT

This VSR documents the results of the validation testing performed on an Ada compiler. Testing was carried out for the following purposes:

- . To attempt to identify any language constructs supported by the compiler that do not conform to the Ada Standard
- . Po attempt to lientify any unsupported language constructs required by the Ada Standard.
 - to determine that the implementation-dependent behaviour is allowed by the Ada Standard

where we this compiler was conducted by NCC index the direction of the NCC index the direction of the NCC index the direction of the Ada Malifectian Organisation (AVO). On-site testing was conducted from the first through 4 June 1987 at Alsys Inc. Waltham MA 02154.

THIS VALIDATION SUMMARY REPORT

Consistent with the national laws of the originating country, the AVO may make full and free public disclosure of this report. In the United States, this is provided in accordance with the "Freedom of Information Act" (5 U.S.C. 552). The results of this validation apply only to the computers, operating systems, and compiler versions identified in this report.

The organisations represented on the signature page of this report do not represent or warrant that all statements set forth in this report are accurate and complete, or that the subject compiler has no nonconformities to the Ada Standard other than those presented. Copies of this report are available to the public from:

Ada Information Clearinghouse Ada Joint Program Office OUSDRE The Pentagon, Rm 3D-139 (Fern Street) Washington DC 20301-3081

or from:

Ada Validation Facility
The National Computing Centre Ltd
Oxford Road
Manchester
M1 7ED
United Kingdom

Questions regarding this report or the validation test results should be directed to the AVF listed above or to:

Ada Validation Organization Institute for Defense Analyses 1801 North Beauregard Alexandria VA 22311

A A REPORT TRENGTES

- Reference Manual for the Ada Programming Language, ANST/MIL-STD-1815A, FEB 183.
- Ada Validation Organization: Policies and Proceduces, MURE Corporation, JUN 1982, 98 83-110601.
- Ada Compiler Validation Capability Implementer's Guide, Soffech, Inc., DEC 1984.

1.4 DEFINITION OF TERMS

CONTROL OF STREET CONTROL OF STREET STREET, ST

The Ada Compiler Validation Capability. A set of programs that evaluates the conformity of a compiler to the Ada language specification, ANSI/MIL-STD-1815A.

Add Standard ANSI/MCL-STD-1815A, February 1983.

Applicanv The agency requesting validation.

AVF The National Computing Centre Ltd. In the context of this report, the AVF is responsible for conducting compiler validations according to established policies and procedures.

AVO The Ada Validation Organization. In the context of this report, the AVO is responsible for setting procedures for compiler validations.

Compiler A processor for the Ada language. In the context of this report, a compiler is any language processor, including cross-compilers, translators, and interpreters.

Failed test A test for which the compiler generates a result that demonstrates nonconformity to the Ada Standard.

Host The computer on which the compiler resides.

Inapplicable A test that uses features of the language that a compiler test is not required to support or may legitimately support in a way other than the one expected by the test.

Perhaps that A test for which a compiler generates the expected result.

The computer (c. which a compiler generates code.

A program that checks a compiler's conformity regarding a particula. feature or features to the Ada Standard. In the context of this report, the term is used to design te a single test, which may comprise one or more files.

A test found to be incorrect and not used to check conformity to test the Ada language specification. A test may be incorrect because it has an invalid test objective, fails to meet its test objective, or contains illegal or erroneous use of the language.

1.5 ACVC TEST CLASSES

edioscopica respesso respendi esperanto presenta esperanto presenta esperanto de presenta esperanto esperanto e

Condormity to the Ada Standard is measured using the ACVC. The ACVC reference both legal and illegal Ada programs structured into six test classes. A, B, C, D, E, and L. The first letter of a test name identifies the class to which it belongs. Class A, C, D, and E tests are executable, and special program units are used to report their results during execution. Class B tests are expected to produce domption errors. Class E tests are expected to produce domption errors.

Class A tests check that legal Ada programs can be successfully compiled and executed. However, no checks are performed during execution to see if the test objective has been met. For example, a Class A test checks that reserved words of another language (other than those already reserved in the Ada language) are not treated as reserved words by an Ada compiler. A Class A test is passed if no errors are detected at compile time and the program executes to produce a PASSED message.

Class B tests check that a compiler detects illegal language usage. Class B tests are not executable. Each test in this class is compiled and the resulting compilation listing is examined to verify that every syntax or semantic error in the test is detected. A Class B test is passed if every illegal construct that it contains is detected by the compiler.

Class C tests check that legal Ada programs can be correctly compiled and executed. Each Class C test is self-checking and produces a PASSED, FAILED, or NOT APPLICABLE message indicating the result when it is executed.

Class D tests check the compilation and execution capabilities of a compiler. Since there are no requirements placed on a compiler by the Ada Standard for some parameters—for example, the number of like this capacity in a compilation or the number of units in a cary a compiler may refuse to compile a Class D test and still be to be a compiler. Therefore, if a Class D test fails to compile the department of the compiler is exceeded, the test is a sinapplicable. If a Class D test compiles successfully, solf—checking and occdudes a CASSED or FAILED message during

Ass E test is self-checking and produces a NOT APPLICABLE, 2N 380, or FAILED message when it is compiled and executed. Mowever, the Ada Standard permits an implementation to reject programs containing some features addressed by Class E tests during co-pilation. Therefore, a Class E test is passed by a compiler if it is compiled successfully and executes to produce a PASSED message, or it it is rejected by the compiler for an allowable reason.

cause A tests check that incomplete or illegal Ada programs involving multiple, separately compiled units are detected and not allowed to execute. Class L tests are compiled separately and execution is authorized. A Class L test passes if it is rejected at link time--that is an autempt to execute the main program must generate an error message before any declarations in the main program or any units referenced by the main program are elaborated.

alleannen sereren perenen errenen sereren errenen kontrolen berenen propertien berenen kontrolen

Two library units, the package REPORT and the procedure CHECK FILE, support the self-checking features of the executable tests. package REPORT provides the mechanism by which executable tests report PASSED, FAILED, or NOT APPLICABLE results. It also provides a set of identity functions used to defeat some compiler optimization allowed by the Ada Standard that would circumvent a test objective. The procedure CHECK FILE is used to check the contents of text files written by some of the Class C tests for chapter 14 of the Ada The operation of these units is checked by a set of Stanlard. executable tests. These tests produce messages that are examined to verify that the units are operating correctly. If these units are not operating correctly, then the validation is not attempted.

INTRODUCTION

The text of the tests in the ACVC follow conventions that are intended to ensure that the tests are reasonably portable without modification. For example, the tests make use of only the basic set of 55 characters, contain lines with a maximum length of 72 characters, use small numeric values, and place features that may not be supported by all implementations in separate tests. However, some tests contain values that require the test to be customized according to applementation specific values for example, an illegal file name. A line values used for this validation is listed in Appendix C.

the mest contently process each of the tests in the suite and make a stell contormity to the Ada Standard either meeting the pass who are given for the test or by showing that the best is able to the implementation. Any test that was determined to an illegal language construct or an erroneous language of its lithdrawn from the NCMC and, therefore, is not used in a compiler. The tests withdrawn at the time of validation are to Appendix D.

SECTION SECTION SECTIONS CONTRACTOR OF SECTION SECTIONS S

CHAPTER 2

CONFIGURATION INFORMATION

- WATTGURATION TESTED

where condition by tem for this validation was tested under the α flowing configuration:

moster: Alsycomp_003, Vs.1

ersion: 1.8

Proceeding action Expiration Date:

F - Computer:

Machine :

Wang PC 280

Operating System:

MS-DOS 3.20

Memory Size:

Machine 1 and Machine 2

Base memory 640K Extended memory 4 MB

Hard disk 33 MB

Parger Computer:

Same as host

Communications Network:

Not applicable

2.2 IMPLEMENTATION CHARACTERISTICS

One of the purposes of validating compilers is to determine the behaviour of a compiler in those areas of the Ada Standard that permit implementations to differ. Class D and E tests specifically check for the implementation differences. However, tests in other classes also characterize an implementation. This compiler is characterized by the following interpretations of the Ada Standard:

f year ies.

The compliant correctly processes compilations containing loop statements he ted to 17 levels, block statements nested to 65 levels, and recursive procedures esparately compiled as subunits nested to 12 levels. It correctly processes a compilation containing 723 variables in the same declarative part. (See tests D55A03A..H (8 tests), D56001B, D64005E..G (3 tests), and D29002K.)

. Universal integer calculations.

An implementation is allowed to reject universal integer calculations having values that exceed SYSTEM.MAX_INT. This implementation does not reject such calculations and processes them correctly. (See tests D4A002A, D4A002B, D4A004A, and D4A004B.)

eredefined types.

This implementation supports the additional predefined types SHORT INTEGER, LONG INTEGER and LONG FLOAT, in the package STANDARD. (See tests B86001C and B86001D.)

Based literals.

and obsessed acceptor securities of the property of the security

An implementation is allowed to reject a based literal with a value exceeding SYSTEM.MAX_INT during compilation, or it may raise NUMERIC_ERROR or CONSTRAINT_ERROR during execution. This implementation raises NUMERIC_ERROR during execution. (See test E24101A.)

Array Types.

An implementation is allowed to raise NUMERIC_ERROR or CONSTRAINT_ERROR for an array having a 'LENGTH that exceeds STANDARD.INTEGER'LAST and/ or SYSTEM.MAX INT.

A packed BOOLEAN array having a 'LENGTH exceeding INTEGER'LAST raises no exception when the array type is declared or when the array objects are declared or illied. (Toe test Chillost.)

packed two-dimensional 30068AN array with more than inTEGER'IAST components raides CONSTRAINT ERROR than the length of a dimension is calculated and exceeds INDERNIAST. The best 052104V.)

null array with one dimension of length greater than ARCEGER LAST may cause NUMERIC_ERROR or CONSTRAINT_ERROR when the declared of assigned. Alternatively, an applementation may accept the declaration. However, lengths must match in array slice assignments. This implementation does not raise NUMERIC_ERROR when the array type is declared. (See test E52103Y.)

In assigning one-dimensional array types, the expression appears to be evaluated in its entirety before CONSTRAINT ERROR is raised when checking whether expression's subtype is compatible with the target's subtype. In assigning two-dimensional array types, the expression does appear to be evaluated in aut its entirety before raised when checking CONSTRAINT ERROR is whether expression's subtype is compatible with the target's subtype. (See test C52013A.)

Discriminated types.

During compilation, an implementation is allowed to either accept or reject an incomplete type with discriminants that is used in an access type definition with a compatible discriminant constraint. This implementation accepts such subtype indications. (See test E38104A.)

on assigning record types with discriminants, the expression appears to be evaluated in its entirety before CONSTRAINT ERROR is raised when checking whether the expression's subtype is compatible with the target's subtype. (See test C52013A.)

. Aggregates.

In the evaluation of a multi-dimensional aggregate, all choices appear to be evaluated before checking against the index type. (See tests C43207A and C43207B.)

In the evaluation of an aggregate containing subaggregates, all choices are not evaluated before being checked for dentical bounds. (See test E4.212B.)

the choices are evaluated before CONSTRAINT ERROR is caused in a bound in a nonnull range of a nonnull aggregate does not colong to an index subtype. (See test E43211B.)

- actions

An implementation may allow the declaration of a parameterless function and an enumeration literal having the same profile in the same immediate scope, or it may reject the function declaration. If it accepts the function declarations, the use of the enumeration literal's identifier denotes the function. This implementation rejects the declarations. (See test E66001D.)

Representation clauses.

The Ada Standard does not require an implementation to support representation clauses. If a representation clause is not supported, then the implementation must reject it. While the operation of representation clauses is not checked by Version 1.8 of the ACVC, they are used in testing other language This implementation accepts 'SIZE for tasks, features. it rejects 'STORAGE SIZE for tasks, 'SMALL clauses; collections. Enumeration representation clauses, including those that specify noncontiguous values, appear to (See tests C55B16A, C87B62A, C87B62B, C87B62C, and supported. BC1002A.)

. Pragmas.

esso escares imperior ballica assessi proporti college, essessi pinanci, essessi proporti essesse.

The pragma INLINE is supported for procedures and functions. (See tests CA3004E and CA3004F.)

. Input/Output.

The package SEQUENTIAL IO can be instantiated with unconstrained array types and record types with discriminants. The package DIRECT IO can be instantiated with unconstrained array types and record types with discriminants without defaults. (See tests AE2101C, AE2101H, CE2201D, CE2201E, and CE2401D.)

CONFIGURATION INFORMATION

An existing text file can be opened in OUT_FILE mode, can be created in OUT_FILE mode, and cannot be created in IN_FILE mode. (See test EE3102C.)

More than one internal file can be associated with each external file for text I/C for reading only. (See tests CE3111A.E (5 tests).)

More than one internal like can be associated with each external falls for sequential I/O for reading only. See tests CE2107A. S (6 tests).)

Fore than one internal file can be associated with each external file for direct T/O for reading only. (See tests = E2107A... (6 tests).)

and external file associated with more than one internal file cannot be deleted. (See test CE2:108.)

Temporary sequential and direct files are given a name. Temporary files given names are not deleted when they are closed. (See tests CE2108A and CE2108C.)

. Generics.

TO SERVICE TO SERVICE AND ASSESSED AND ASSESSED IN THE SERVICE ASSESSED ASSESSED ASSESSED ASSESSED.

Generic subprogram declarations and bodies can be compiled in separate compilations. (See tests CA2009F.)

Seneric package declarations and bodies can be compiled in separate compilations. (See tests CA2009C and BC3205D.)

CHAPTER 3

TEST INFORMATION

THE PROOF RESULTS

The AVF determined that 196 tests were inapplicable to this lengthmentation, and that the 2184 applicable tests were passed by the implementation.

The AVF concludes that the testing results demonstrate acceptable conformity to the Ada Standard.

J. J. SAMMARY OF TEST RESULTS BY CLASS

$\mathcal{A}U\mathcal{B}$		TEST	CLASS				TOTAL
	A	В	c	D	E	L	
Passed	69	865	1178	13	13	46	2184
Failed	0	0	Ü	0	0	0	0
Inapplicable	0	2	190	4	0	0	196
Withdrawn	0	7	12	0	0	0	19
TOTAL	69	874	1380	17	13	46	2399

3.3 SUMMARY OF TEST RESULTS BY CHAPTER

RESULT					C	HAP	rer						
	2	3	4	5	6	7	8	9	10	11	12	14	TOTAL
°3⇔ sed	102	253	334	343	ι61	97	136	261	129	32	218	518	2184
Marriad	ŋ	0	0	٠,	0	0	υ	0	0	O	0	U	۴.
framplicab).	14	72.	36	4	0	0	3	ι	1.	0	0	L5	196
d - Atawn	O	5	5	o	o	3	Į	Ż	ů,	0	1	0	19
1. 191	116	330	425	247	161	98	140	264	134	39	219	233	2399

WITHORAWN TESTS

The following 19 tests were withdrawn from ACVC Version 1.8 at the time of this validation:

C32114A	C41404A	B74101B
B33203C	B45116A	C87B50A
C04918A	C48008A	C92005A
C35904A	B49006A	C940ACA
33740 1A	B4A010C	CA3005AD (4 tests)
		BC3204C

See Appendix D for the reason that each of these tests was

3.5 INAPPLICABLE TESTS

dome tests do not apply to all compilers because they make use of features that a compiler is not required by the Ada Standard to supposite. Others may depend on the result of another test that is either inapplicable or withdrawn. For this validation attempt, 196 tests were inapplicable for the reasons indicated:

- . C34001F and C35702A use SHORT_FLOAT which is not supported by this compiler.
- . D55A03E..H (4 tests) because the compiler only processes compilations containing loop statements nested to 17 levels.

- . B86001D requires a predefined numeric type other than those defined by the Ada language in package STANDARD. There is no such type for this implementation.
- . C86001F redefines package SYSTEM, but TEXT_IO is made obsolete by this new definition in this implementation and the test cannot be executed since the package REPORT is dependent on the package TEXT_IO.
 - a uses the 'STORAGE SIVE clause to specify the mollection size for an access type which is not supported by this compiler. The STORAGE SIZE clause is rejected during compilation.
 - controllers implementations for which the smallest and largest values in type DURATION are different from the smallest and largest values in DURATION's base type. This is not the case for this implementation.
- BA2001E requires that duplicate names of subunits with a common ancestor be detected at compilation time. This compiler correctly detects the error at link time and the AVO rules that such behaviour is acceptable.
- CE21078..E (4 tests), CE2110B, CE2111D, CE2111H, CE3111B..E (4 tests), CE3114B, and CE3115A are inapplicable because multiple internal files cannot be associated with the same external file objective file is open for writing. The proper exception is raised at tiple access is attempted.
 - (821020) because mode (N_FILE is not supported for SEQUENTIAL_IO.
 - 2221021 because mode IN_FILE is not supported for DIRECT_IO.
- The following 170 tests make use of floating-point precision that exceeds the maximum of 15 supported by the implementation:
 - C24113L..Y (14 tests)
 C35705L..Y (14 tests)
 C35706L..Y (14 tests)
 C35707L..Y (14 tests)
 C35708L..Y (14 tests)
 C35802L..Y (14 tests)
 C45241L..Y (14 tests)
 C45241L..Y (14 tests)
 C45421L..Y (14 tests)
 C45421L..Y (14 tests)
 C45421L..Y (14 tests)
 C45621L..Z (15 tests)
 C45621L..Z (15 tests)

ged between the property of th

3.6 SPLIT TESTS

If one or more errors do not appear to have been detected in a Class B test because of compiler error recovery, then the test is split into a set of smaller tests that contain the undetected errors. These splits are then compiled and examined. The splitting process continues until all errors are detected by the compiler or until there is exactly one per split. Any Class A, Class C, or Class E test that cannot be easily and executed because of its size is split into a second obsets that can be processed.

System one required for 14 Class B tests.

82 600 5A	B33001A	B37004A
843201D	B45102A	Bo (012A
3620018	B62001C	874401F
374401R	B91004A	895069A
B95069B	8C3205C	

ADDITIONAL TESTING INFORMATION

. / : Prevalidation

process assessed sociality isolation population

Prior to validation, a set of test results for ACVC Version 1.8 produced by AlsyCOMP_003 was submitted to the AVF by the applicant for test end Analysis of these results demonstrated that the compiler successfully passed all applicable tests, and the compiler exhibited the expected behaviour on all inapplicable tests.

3 'est Method

Testing of AlsyCOMP_003 using ACVC Version 1.8 was conducted on-site by a validation team from the AVF. The configuration consisted of two Wang PC 280s operating under MS-DOS 3.20 which were both host and target.

A magnetic tape containing all tests was taken on site by the validation team for processing. The magnetic tape contained tests that make use of implementation-specific values were customized being written to the magnetic tape. Tests requiring splits during the prevalidation testing were not included in their split form on the magnetic tape.

The contents of the magnetic tape were loaded first onto a VAX 750 computer, where the required splits were performed. The processed source files were then transferred to an IBM PC AT computer via an Ethernet connection. Two Wang PC 280 machines were used to process the validation suite. Source files were then transferred to the host machines via floppy disks.

After the test files were loaded to disk, the full set of tests was compiled and linked on the Wang PC 280s and all executable tests were run. Results were written to floppy disk, moved to an IBM PC and then transferred to a VAX 750 via Ethernet. The results were then printed from the VAX 750.

The compiler was tested using command scripts provided by Alsys Inc. and eviewed by the validation team. The following options were in for testing:

potion Effect
This option allows insertion of code for subprograms inline and must be set for the pragma INLIME to be operative.

ere compiled, linked and executed (as appropriate) using a look lost computer and a single target computer. Test output, compiled ion listings, and job logs were captured on magnetic tape and archived at AVF. The listings examined on-site by the validation lead were also archived.

3.7.3 WEST SITE

Secret Variable Consisted Properties and Secretion and Secretion Consisted and Consisted Secretion of Consisted Consisted Secretion of Consisted Consisted Secretion of Consisted Consiste

The validation team arrived at Alsys Inc. Waltham MA 02154 on 1 June 1987 and departed after testing was completed on 4 June 1987.

APPENDIX A

COMPLIANCE STATEMENT

Alsys inc has submitted the following compliance statement concerning the Aleycomp 00%.

COMPLIANCE STATEMENT

	Complian	nce Statement	
Base $C(mC_F)$ rotion:			
i er:	AlsyCO or 1003 Wars	sion 3.1	
e est suite:	Ada* Compiler Valids	nto. Capability, Version 1.8	
146 computer:			
	Sischine:	Wang (%) 280	
	Operating System:	MS-DOS 3,20	
tacker Computer:	Same as host		
Communications Ne	twork:	Not applicable	
Alsys Inc has made no delifalsys Inchagages to the pub	lic disclosure of this rep		da Joine Program Offic
Alsys Inc Arra Avakian Vice President of Engineer		Date:	June 4, 1987

*Ada is registered trademark of the United States Government

(Ada Joint Program Office)

APPENDIX B

APPENDIX F OF THE Ada STANDARD

oniy allowed implementation dependencies orrasboud was taken to n-dependent pragm i, to certain machine dependent of welling, as mentioned in chapter (3 of MILESTD-1818), restrictions on representation :11owed classes. -materion-dependent characteristics of the AlsyCoMP 003, V3.1 are in the following sections which discuss topics one through stated in Appendix F of the Ada Language Reference Manual (Advay If AD-16 LoA). {Implementation-specific portions l'ANDARD are also included in this appendix. The space first for of the package STANDARD is also included in this or endiv.

Package STANDARD is

```
* 1 °
```

and because assessed basearer personal personal interestal property becauses accuracy basearer besones

Type INTEGER is -32 768 .. 32 767;

Type SHORT INTEGER is range -128 .. 127;

Type FONG INTEGER is range -2 147 483 648 .. 2 147 483 647;

type + LOAT is digits 6 range -2#1.111_1111_1111_1111_1111#E+127 .. 2#1.111_1111_1111_1111_1111#E+127;

type DURATION is delta 0.001 range -2 097 152.0 .. 2 097 151.999;

end STANDARD;

Alsys PC At Ada Compiler

APPENDIX F

Implementation - Dependent Characteristics

Version 3.1

Alsys S.A. 29. Avenue de Versailles 78170 La Celle St. Cloud, France

Alsys, Inc. 1432 Main Street Waltham, MA 02154, U.S.A.

Alsys Ltd.
Partridge House, Newtown Road
Henley-on-Thames,
Oxon RG9 1EN, England

^{*} Ada is a registered trademark of the U.S. Government, Ada Joint Program Office

Copyright 1987 by Alsys

All rights reserved. No part of this document may be reproduced in any form or by any means without permission in writing from Alsys.

Printed: June 1987

beneral encourage reassesses, recommon appealable encountries are resistion of the contraction of the contra

Alsys reserves the right to make changes in specifications and other information contained in this publication without prior notice. Consult Alsys to determine whether such changes have been made.

^{*} Ada is a registered trademark of the U.S. Government, Ada Joint Program Office

TABLE OF CONTENES

AP	PENDIX F	1
Ť	Implementation-Dependent Pragmas	2
	1.1 INTERFACE 1.2 INTERFACE NAME 1.3 Other Pragmas	2 2 3
2	Implementation - Dependent Attributes	4
3	Specification of the package SYSTEM	4
4	Restrictions on Representation Clauses	7
5	Conventions for Implementation-Generated Names	s 7
6	Address Clauses	8
7	Restrictions on Unchecked Conversions	8
8	Input-Output Packages	8
	8.1 Correspondence between External Files and DOS Files 8.2 Error Handling 8.3 The FORM Parameter 8.4 Sequential Files 8.5 Direct Files 8.6 Text Files 8.7 Access Protection of External Files 8.8 The Need to Close a File Explicitly 8.9 Limitation on the procedure RESET	8 9 9 9 9 10 10
	8.9 Limitation on the procedure RESET 8.10 Sharing of External Files and Tasking Issues	

Table of Contents

9	Cha	racteristics of Numeric Types	11
	9.1 9.2 9.3	Integer Types Ficating Point Type Attributes Attributes of Type DURATION	11 12 12
10	Oth	er Implementation-Dependent Characteristics	13
	0.1	Use of the Ficating (Fig. Coprocessor (808") 80087)	13
	10.2	Champteristics of the Willip	13
	10.3	Characteristics of Taska	13
	10.4	Definition of a Main Subprogram	14
	10.5	Octoring of Campilation Units	14
, i	Limi	tations	15
	11.1	Compiler Limitations	15
	(1.2	Hardware Related Limitations	15
	11.3	Runtime Executive Limitations	15
IND	FΧ		17

APPENDIX F

Implementation - Dependent Characteristics

The appendix summarizes the implementation-dependent characteristics of the Alsys PC Ada Compiler. This appendix is a required part of the Reference Manual for the two Programming Language (called the RM in this appendix).

the sections of this appendix are as follows:

STORY OF STANDED STAND

- 1. The form, allowed places, and effect of every implementation-dependent pragma.
- 2. The name and the type of every implementation-dependent attribute.
- 3. The specification of the package SYSTEM.
- 4. The list of all restrictions on representation clauses.
- 5. The conventions used for any implementation-generated name denoting implementation-dependent components.
- 6. The interpretation of expressions that appear in address clauses, including those for interrupts.
- 7. Any restrictions on unchecked conversions.
- 8. Any implementation-dependent characteristics of the input-output packages.
- 9 Characteristics of numeric types.
- 10. Other implementation-dependent characteristics.
- 11. Compiler limitations.

The name Alsys Runtime Executive Programs or simply Runtime Executive refers to the runtime library routines provided for all Ada programs. These routines implement the Ada heap, exceptions, tasking control, and other utility functions.

General systems programming notes are given in another document, the Application Developer's Guide (for example, parameter passing conventions needed for interface with assembly routines).

1 Implementation-Dependent Pragmas

Ada programs can interface with subprograms written in Assembler and other languages through the use of the predefined pragma INTERFACE and the implementation-defined pragma INTERFACE NAME.

ANTERFACE

Pragma INTERFACE specifies the name of an interfaced subprogram and the name of the programming language for which parameter passing conventions will be generated. Tragma INTERFACE takes the form specified in the RM:

oragina iNTERFACE (language_name, subprogram_name):

vhere

- a language name is ASSEMBLER or ADA.
- subprogram_name is the name used within the Ada program to refer to the interfaced subprogram.

The only two language names accepted by pragma INTERFACE are ASSEMBLER and ADA. The full implementation requirements for writing pragma INTERFACE subprograms are described in the Application Developer's Guide.

The language name used in the pragma INTERFACE does not have to have any relationship to the language actually used to write the interfaced subprogram. It is used only to tell the Compiler how to generate subprogram calls; that is, what kind of parameter passing techniques to use. The programmer can interface Ada programs with subroutines written in any other (compiled) language by understanding the mechanisms used for parameter passing by the Alsys PC AT Ada Compiler and the corresponding mechanisms of the chosen external language.

1.2 INTERFACE NAME

Pragma INTERFACE NAME associates the name of the interfaced subprogram with the external name of the interfaced subprogram. If pragma INTERFACE_NAME is not used, then the two names are assumed to be identical. This pragma takes the form:

pragma INTERFACE_NAME (subprogram_name, string_literal);

where,

- subprogram_name is the name used within the Ada program to refer to the interfaced subprogram.
- string_literal is the name by which the interfaced subprogram is referred to at link time.

The pragma INTERFACE_NAME is used to identify routines in other languages that are not named with legal Ada identifiers. Ada identifiers can only contain letters, digits, or underscores, whereas the DOS Linker allows external names to contain other characters, for example, the dollar sign (\$) or commercial at sign (@). These characters are the specified in the string literal argument of the gragma INTERFACE NAME

the pragma INTERFACE NAME is allowed at the same places of an Ada program as the pragma INTERFACE. (Location restrictions can be found in section 13.9 of the RM.) However, the pragma INTERFACE_NAME must always occur after the pragma INTERFACE declaration for the interfaced subprogram.

The string literal of the pragma INTERFACE_NAME is passed through unchanged to The DOS object (ile. (The DOS tools usually ignore the case of external identifiers. Towever recent versions of the PLINK86 and Microsoft Linkers have options to treat external identifiers in a case-sensitive manner.) The maximum length of the string treral is 40 characters. This limit is not noted by the Compiler, but is truncated by the Binder to meet the Intel object module format standard. Certain DOS tools have smaller limits. (For example, the IBM Macro Assembler limits external identifiers to 31 characters.)

The Runtime Executive contains several external identifiers. All such identifiers begin with either the string "ADAA" or the string "ADAA". Accordingly, names prefixed by "ADAA" or "ADAAA" should be avoided by the user.

Example

processes assessment parameter decreases respected antipical despected between

```
package SAMPLE_DATA is
  function SAMPLE_DAVICE (X: INTEGER) return INTEGER;
  function PROCESS_SAMPLE (X: INTEGER) return INTEGER;
private
  pragma INTERFACE (ASSEMBLER, SAMPLE_DEVICE);
  pragma INTERFACE (ADA, PROCESS_SAMPLE);
  pragma INTERFACE_NAME (SAMPLE_DEVICE, "DEVICSGET_SAMPLE");
end SAMPLE_DATA;
```

1.3 Other Pragmas

Pragma PRIORITY is accepted with the range of priorities running from 1 to 10 (see the definition of the predefined package SYSTEM in Section 3). Undefined priority (no pragma PRIORITY) is treated as though it were less than any defined priority value.

In addition to pragma SUPPRESS, it is possible to suppress all checks in a given compilation by the use of the Compiler option CHECKS.

2 Implementation-Dependent Attributes

PTS_ARRAY

For a prefix P that denotes any type or suctype, this attribute yields the value TRUE if P in an array type or an array subtype; otherwise, it yields the value FALSE.

3 Syncification of the package SYSTEM

```
Williage SYSTEM is
             * (1) Required Definitions.
      type NAME is (1_80x86);
      SYSTEM_NAME : constant NAME := 1_80x86;
      STORAGE_UNIT : constant := 8;
      MEMORY_SIZE : constant := 640 * 1024;
      -- System-Dependent Named Numbers:
      MIN_INT
                  : constant := -(2 **31);
      MAX_INT
                  : constant := 2**31 - 1;
      MAX_DIGITS : constant := 15;
      MAX_MANTISSA : constant := 31;
      FINE DELTA : constant := 2#1.0#E-31;
      -- For the high-resolution timer, the clock resolution is
      -- 1.0 / 1024.0.
                  : constant := 1.0 / 18.2;
      · Other System-Dependent Declarations:
      subtype PRICRITY is INTEGER range 1 .. 10;
      -- Though declared here as access to a STRING in fact this
      ·· can be used anywhere an ADDRESS is required by Ada.
      -- The type ADDRESS is, in fact, implemented as an
      -- 8086/80286 segment:offset pair.
      type ADDRESS is access STRING;
      NULL ADDRESS: constant ADDRESS := null;
```

```
*******
       * (2) MACHINE TYPE CONVERSIONS *
\sim 1f the word / double-word operations below are used on
-- ADDRESS, then MSW yields the segment and LSW yield, the
· offset.
-- In the operations below, a BYTE_TYPE is any simple type
· implemented on 8-bits (for example, SHORT_INTEGER), a WORD_TYPE is
-- any simple type implemented on 16 bits (for example, IMTEGER), and
· a DOUBLE WORD TYPE is any simple type implemented on
- 12 bits (for example, LONG_INTEGER, FLOAT ADDRESS).
 . Byta <==> Word conversions:
-- Get the most significant byte:
generic
      type BYTE_TYPE is private;
      type WORD_TYPE is private;
function MSB (W: WORD_TYPE) return BYTE_TYPE;
-- Get the least significant byte:
generic
      type BYTE_TYPE is private;
      type WCRD_TYPE is private;
function LSB (W: WORD_TYPE) return BYTE_TYPE;
·- Compose a word from two bytes:
generic
      type BYTE_TYPE is private;
      type WORD_TYPE is private;
function WORD (MSB, LSB: BYTE_TYPE) return WORD_TYPE;
-- Word <==> Double-Word conversions:
-- Get the most significant word:
      type WCRD_TYPE is private;
      type DOUBLE_WORD_TYPE is private;
function MSW (W: DOUBLE_WORD_TYPE) return WORD_TYPE;
-- Get the least significant word:
generic
      type WORD_TYPE is private;
      type DOUBLE_WORD_TYPE is private;
function LSW(W: DOUBLE_WORD_TYPE) return WORD_TYPE;
```

5

CONTRACTOR CONTRACTOR

```
· · Compose a DATA double word from two words.
       generic
             type WORD_TYPE is private;
             · The following type must be a data type
             - (for example, LONG INTEGER):
             +.pe DATA_DOUBLE_WORD is private;
       function houses word (MSW, USW: WORD TYPE)
                           return DATA_DOUBLE_MORD;
       -- Compose a REFERENCE double word from two words.
       generic
             type WORD_TYPE is private;
             -- The following type must be a reference type
             ·· (for example, access or ADDRESS):
             type REF_DOUBLE_WORD is private;
       function PEFERENCE (SEGMENT, OFFSET: YORD_TYPE)
                           return REF_DOUBLE_WORD;
              * (3) OPERATIONS ON ADDRESS *
       ·- You can get an address via 'ADDRESS attribute or by
       ·· instantiating the function REFERENCE, above, with
       · appropriate types.
        - Some addresses are used by the Compiler. For example,
       ·· the display is located at the low end of the DS segment,
       -- and addresses SS:0 through SS:128 hold the task control
       · · block and other information. Writing into these areas
       ·· will have unpredictable results.
       \cdot \cdot \cdot In real mode, the memory for DOS itself, including all the
      ·· interrupt vectors is also unprotected. Thus, any user of
       -- ASSIGN_TO_ADDRESS must be extremely careful.
      -- Note that no operations are defined to get the values of
       -- the segment registers, but if it is necessary an
      · interfaced function can be written.
             type OBJECT is private;
      function FETCH_FROM_ADDRESS (FROM: ADDRESS) return OBJECT;
      generic
             type OBJECT is private;
      procedure ASSIGN_TO_ADDRESS (OBJ: OBJECT; TO: ADDRESS);
end SYSTEM;
```

SEFERI LEGISSON OF SERVING THE SERVING CONTINUES. DESIGNATION SERVING CONSISTS CONTINUES.

4 Restrictions on Representation Clauses

The facilities covered in Chapter 13 of the RM are provided, except for the following features:

- · Address clauses are not implemented.
- There is no bit implementation for any of the representation clauses.
- The Record Clause is not allowed for a derived record type.
- Change of representation for RECORD type is not implemented.
- The Enumeration Clause is not allowed if there is a range constraint on the parent subtype.
- For the length clause:
 - Size specification: T'SIZE is not implemented for types declared in a generic unit.
 - Specification of storage for a task activation: TSTORAGE_SIZE is not implemented when T is a task type.
 - Specification of small for a fixed point type: TSMALL is restricted to a power of 2, and the absolute value of the exponent must be less than 31.
- Machine code insertions are not implemented; use instead pragma INTER-FACE to ASSEMBLER to write assembly language routines.

5 Conventions for Implementation-Generated Names

The Alsys PC AT Ada Compiler may add fields to record objects and have descriptors in memory for record or array objects. These fields are not accessible to the user through any implementation-generated name or attribute.

The following predefined packages are reserved to Alsys and cannot be recompiled in Version 3.1:

ALSYS_ADA_RUNTIME
ALSYS_BASIC_IO
ALSYS_BASIC_DIRECT_IO
ALSYS_BASIC_SEQUENTIAL_IO

6 Address Clauses

This version of the Alsys PC AF Ada Compiler does not support address clauses. Support is provided for Ada interrupt entries; please see Chapter 6 of the Application Description of the Alsys PC AF Ada Compiler does not support address clauses. Support is provided for Ada interrupt entries; please see Chapter 6 of the Application Description of the Application Descriptio

Restrictions on Unchecked Conversions

sensibled conversions are allowed between any types. It is the programmer's re-

💲 — Input-Output Packages

The XM defines the predefined input-output packages SEQUENTIAL_IO, DIRECT_IO, and IEXT_IO, and describes how to use the facilities available within these packages. The RM also defines the package IO_EXCEPTIONS, which specifies the exceptions that can be raised by the predefined input-output packages.

In addition the RM outlines the package LOW_LEVEL_IO, which is concerned with low-level machine-dependent input-output, such as would possibly be used to write device drivers or access device registers. LOW_LEVEL_IO has not been implemented. The use of interfaced subprograms is recommended as an alternative.

8.1 Correspondence between External Files and DOS Files

Add input-output is defined in terms of external files. Data is read from and written to external files. Each external file is implemented as a standard DOS file, including the use of STANDARD_INPUT and STANDARD_OUTPUT.

The name of an external file can be either

- the null string
- a DOS filename
- a DOS special file or device name (for example, CON and PRN)

If the name is a null string, the associated external file is a temporary file and will cease to exist when the program is terminated. The file will be placed in the current directory and its name will be chosen by DOS.

If the name is a DOS filename, the filename will be interpreted according to standard DOS conventions (that is, relative to the current directory). The exception NAME_ERROR is raised if the name part of the filename has more than 8 characters or if the extension part has more than 3 characters.

If an existing DOS file is specified to the CREATE procedure, the contents of the file will be deleted before writing to the file.

CONTINUE CONSTRUE CONTRACTO RESISTANT CRICIARIO REFERENCE CONTRACTO DESCRIPTO DE

If a non-existing directory is specified in a file path name to CREATE, the directory will not be created, and the exception NAME_ERROR is raised.

Tor Handling

DO LONS are translated into Ada exceptions, as defined in the RM by package 30 EMCEPTIONS. In particular, DEVICE_ERROR is raised in cases of drive not ready, whom media, disk full or hardware errors on the disk (such as read or write rauli).

8.3 FORM Parameter

The only accepted value of the FORM parameter of the various CREATE and OPEN procedures is the null string. The functions FORM of TEXT_IO and (any instantiation of) DIRECT_IO and SEQUENTIAL_IO always return the null string. USE_ERROR is raised if a non-null FORM string is passed to any CREATE or OPEN procedure.

8.4 Sequential Files

For sequential access the file is viewed as a sequence of values that are transferred in the order of their appearance (as produced by the program or run-time environment). This is sometimes called a *stream* file in other operating systems. Each object in a sequential file has the same binary representation as the Ada object in the executable program.

8.5 Direct Files

For direct access the file is viewed as a set of elements occupying consecutive positions in a linear order. The position of an element in a direct file is specified by its index, which is an integer of subtype POSITIVE COUNT.

DIRECT_IO only allows input-output for constrained types. If DIRECT_IO is instantiated for an unconstrained type, all calls to CREATE or OPEN will raise USE_ERROR. Each object in a direct file will have the same binary representation as the Ada object in the executable program. All elements within the file will have the same length.

8.6 Text Files

Text files are used for the input and output of information in ASCII character form. Each text file is a sequence of characters grouped into lines, and lines are grouped into a sequence of pages.

All text file column numbers, line numbers, and page numbers are values of the subtype POSITIVE_COUNT.

Note that due to the definitions of line terminator, page terminator, and file terminator in the RM, and the method used to mark the end of file under DOS, some ASCII files do not represent well-formed TEXT IO files.

A text file is buffered by the Runtime Executive unless

- it trames a device (for example, CON or PRN).
- It is STANDARD_INPUT or STANDARD_OUTPUT and has not been redirected.

If not redirected, prompts written to STANDARD OUTPUT with the procedure PUT will appear before (or when) a GET (or GET [INE]) occurs.

The forctions END OF PAGE and END OF FILE always return FALSE when the life is a device, which includes the use of the file CON, and STANDARD_INPUt when it is not redirected. Programs which would like to check for end of file when the file may be a device should handle the exception END_ERROR instead, as in the following example:

Example

END_ERROR is raised for STANDARD_INPUT when ^Z (ASCII.SUB) is entered at the keyboard.

8.7 Access Protection of External Files

Standard DOS does not provide any access protection for external files. In a network environment, access protection is dependent on the network operating system.

8.8 The Need to Close a File Explicitly

The Runtime Executive will flush all buffers and close all open files when the program is terminated, either normally or through some exception.

CONTRACTOR OF STANDARD CONTRACTOR OF THE ASSESSMENT OF THE CONTRACTOR OF THE CONTRAC

However, the RM does not define what happens when a program terminates without closing all the opened files. Thus a program which depends on this feature of the satisfactories might have problems when ported to another system.

39 Limitation on the procedure RESET

THE BESSESSE BESTELL PROPERTY SERVING PROPERTY.

An internal file opened for input cannot be RESET for output. However, an internal file opened for output can be RESET for input, and can subsequently be RESET back to output.

3.40 Sharing of External Files and Tasking Issues

leveral internal files can be associated with the same external file only if all the internal files are opened with mode fN_MODE. However, if a file is opened with mode CUT_MODE and then changed to IN_MODE with the RESET procedure, it cannot be second.

Care should be taken when performing multiple input-output operations on an external file during tasking because the order of calls to the I/O primitives is unpredictable. For example, two strings output by TEXT_IO.PUT_LINE in two different tasks may appear in the output file with interleaved characters. Synchronization of I/O in cases such as this is the user's responsibility.

The TEXT_IO files STANDARD_INPUT and STANDARD_OUTPUT are shared by all tasks of an Ada program.

If fEXF_IO.STANDARD_INPUT is not redirected, it will not block a program on input. All tasks not waiting for input will continue running.

9 Characteristics of Numeric Types

9.1 Integer Types

The ranges of values for integer types declared in package STANDARD are as follows:

SHOR)_INTEGER -128 .. 127 -- 2**7 - 1

INTEGER -32768 .. 32767 -- 2**15 - 1

LONG_INTEGER -2147483648 .. 2147483647 -- 2**31 - 1

For the packages DIRECT_IO and TEXT_IO, the range of values for types COUNT and POSITIVE COUNT are as follows:

COUNT 0 .. 2147483647 .. 2**31 · 1

POSITIVE_COUNT 1 .. 2147483647 .. 2**31 · 1

For the package TEXT_IO, the range of values for the type FIELD is as follows:

FIELD

0 .. 255

2**8 - 1

32 Sloating Point Type Attributes

	FLONE	CONGUELMAT
016115	5	15
MANTISS ⁴	21	51
EMAX.	84	204
EPSILON	9.536748-07	8.88178E-16
LARGE	1.93428E+25	2.57110E+61
SAFE_EMAX	125	1021
SAFE_SMALL	1.17549E-38	2.22507E-308
SAFE_LARGE	4.25353E+37	2.24712E+307
FIRST	-3.40282E+38	-1.79769E+308
LAST	3.40282E+38	1.79769E+308
MACHIRE_RADIX	2	2
MACHINE_MANTISSA	24	53
MACHINE_EMAX	128	1024
MACHINE_SHIN	-125	-1021
MACHINE_ROUNDS	true	true
MACHINE_OVERFLOWS	false	false
SIZE	32	64

9.3 Attributes of Type DURATION

DURATION DELTA

0.001

DURATION'SMALL

0.0009765625 (= 2**(-10))

DURATION'FIRST

-2097152.0

aca hanna e seesee acades posses passes espesse espesse espesse accesse espesses espesses espesses espesses es

DURATION LAST 2097151.999

DURATION LARGE

same as DURATION'LAST

10 Other Implementation-Dependent Characteristics

10.1 Use of the Floating-Point Coprocessor (8087 or 40287)

The Alsys PC AT Ada Compiler generates instructions to use the floating point copropessor for all floating point operations (but, of course, not for operations involving only mersal_real).

Coating point coprocessor, 8087 or 80287, is required for the execution of programs that use arithmetic on floating point values. The coprocessor is needed if the FLOAT_IO or FIXED_IO packages of TEXT_IO are used.

For the AT, the Runtime Executive will detect the absence of the floating point coprocessor if it is required by a program and will raise NUMERIC_ERROR. The user is warned to not run a floating point program on an XT without an 8087. The result of doing so is that the XT will wait indefinitely until a power on reset.

10.2 Characteristics of the Heap

UNCHECKED_DEALLOCATION is implemented for all Ada access objects except access objects to tasks. Use of UNCHECKED_DEALLOCATION on a task object will lead to unpredictable results.

There is no management of collections (via the 'STORAGE_SIZE representation clause on access type) although small objects are managed more efficiently than normal heap objects.

All objects whose visibility is linked to a subprogram, task body, or block have their storage reclaimed at exit.

The maximum size of the heap is limited only by available memory. User programs making use of the Protected Mode feature have a heap which includes all available memory below the 640K limit plus the size of the heap file which is in the extended memory RAM disk.

All objects created by allocators go into the heap. Also, portions of the Runtime Executive representation of task objects, including the task stacks, are allocated in the heap.

10.3 Characteristics of Tasks

The default task stack size is 1K bytes (32K bytes for the environment task), but by using the Binder option STACK. TASK or the SETOPT program, the size for all task stacks in a program may be set to a size from 1K bytes to 64K bytes.

Normal priority rules are followed for preemption, where PRIORITY values are in the range 1...10. A task with undefined priority (no pragma PRIORITY) is considered to be lower than priority 1.

The minimum timeable delay is

- 10/18.2 seconds on a PC XT, or
- 1.0/1024.0 seconds on a PC AT when the high-resolution timer is in effect through the TIMER = FAST option of the BIND command.

the chaximum aumber of active tasks is restricted only by memory usage.

It is opened of a rendezvous executes the accept body code in its own stack. Rendezvous with an empty accept body (for synchronization) does not cause a context switch

The main program waits for completion of all tasks dependent upon library packages before terminating.

Abnormal completion of an aborted task takes place immediately, except when the abnormal task is the caller of an entry that is engaged in a rendezvous, or if it is in the process of activating some tasks. Any such task becomes abnormally completed as soon as the state in question is exited.

The message

THE PARTY OF THE PROPERTY OF THE PARTY OF TH

GLOBAL BLOCKING SITUATION DETECTED

is printed to STANDARD_OUTPUT when the Runtime Executive detects that no further progress is possible for any task in the program. The execution of the program is then abandoned.

10.4 Definition of a Main Subprogram

A liberry unit can be used as a main subprogram if and only if it is a procedure that is not generic and that has no formal parameters.

10.5 Ordering of Compilation Units

The Alsys PC AT Ada Compiler imposes no additional ordering constraints on compilations beyond those required by the language. However, if a generic unit is instantiated during a compilation, its body must be compiled prior to the completion of that compilation.

11 Limitations

11.1 Compiler Limitations

- The maximum identifier length is 255 characters.
- The maximum line length is 255 characters.
- The maximum number of unique identifiers per compilation unit is 2500.
- The maximum number of compilation units in a library is 1000.

11.2 Hardware Related Limitations

- The maximum size of the generated code for a single compilation unit is 65535 bytes.
- The maximum size of a single array or record object is 65522 bytes. The maximum size of a static record is 4096 bytes.
- The maximum size of a single stack frame is 32766 bytes, including the data for inner package subunits "unnested" to the parent frame.
- The maximum amount of data in the global data area is 65535 bytes, including compiler generated data that goes into the GDA (about 8 bytes per compilation unit plus 4 bytes per externally visible subprogram).
- The maximum amount of data in the heap is limited only by available memory.

11.3 Runtime Executive Limitations

Task stacks are only allocated in low memory (that is, below 640k).

INDEX

Abreenal completion 14	Drive not ready 9
• Aborted task 14	DURATION'DELTA 12
Access protection 10	DURATION'FIRST 12
Address clauses 7, 8	DURATION'LARGE 13
Allocators 13	DURATION'LAST 13
Application Developer's Guide 2, 8	DURATION'SMALL 12
Array objects 7	
Array subtype 4	80287 13
Array type 4	8087 13
SSIGN_TO_ADDRESS 6	EMAX 12
Attributes of type DURATION 12	Empty accept body 14
Atterpates of type DORATION 12	END ERROR 10
Binde: 13	END OF FILE 10
	END_OF_PAGE 10 END_OF_PAGE 10
Bit representation clauses 7	
Buffered files 10	Enumeration Clause 7
Buffers	EPSILON 12
flushing 10	Errors disk full 9
Character 2	
Change of representation 7	drive not ready 9 hardware 9
Characteristics of tasks 13	
Collection management 13	unknown media 9
Column numbers 9	FETOU FROM A PRRESS (
Compiler limitations 15	FETCH_FROM_ADDRESS 6
maximum identifier length 15	FIELD 12
maximum line length 15	File closing
maximum number of compilation	explicit 10
units 15	File names 8
maximum number of unique	File terminator 10
identifiers 15	FIRST 12
Constrained types	Fixed point type 7
I/O on 9	FIXED_IO 13
Control Z 10	FLOAT_IO 13
COUNT II	Floating point coprocessor 13
CREATE 8, 9	Floating point operations 13
	Floating point type attributes 12
Device name 8	FORM parameter 9
DEVICE_ERROR 9	
DIGITS 12	GET 10
Direct files 9	GET_LINE 10
DIRECT_10 8, 9, 11	GLOBAL BLOCKING SITUATION
Disk full 9	DETECTED 14
DOS conventions 8	
DOS errors 9	Hardware errors 9
DOS files 8	Hardware limitations
DOS Linker 3	maximum amount of data in the
DOS special file 8	global data area 15

17

, -	maximum data in the heap 15 maximum size of a single array or record object 15 maximum size of a single stack frame 15 maximum size of the generated rode 15 Handa to related limitations 15 Heap 13	Maximum number of unique identifiers 15 Maximum size of a single array or record object 15 Maximum size of a single stack frame 15 Maximum size of the generated code 15 Microsoft Linkers 3 Minimum timeable delay 14
	I/O synchronization 11	NAME_ERROR 8, 9
	IBM Macro Assembler 3	Non-blocking I/O 11
	Implementation generated names 7	Number of active tasks 14
	N MODE II	NUMERIC ERROR 13
	WITEGER II	
	Integer types 11	OPEN 9
	intel object module format 3	Ordering of compilation units 14
	NTERFACE 2, 3	OUT_MODE II
	INTERFACE_NAME 2, 3	
	Interfaced subprograms 8	P'IS_ARRAY 4
	Interleaved characters 11	Page numbers 9
	IO_EXCEPTIONS 8, 9	Page terminator 10
		Parameter passing 1
	LARGE 12	PLINK86 3
	LAST 12	POSITIVE_COUNT 9, 11
	Legal file names 8	Pragma INTERFACE 2, 3 Pragma INTERFACE_NAME 2, 3
	Length clause 7 Library unit 14	Pragma PRIORITY 3, 14
	Limitations 15	Pragma SUPPRESS 3
	Line numbers 9	Predefined packages 7
	Line terminator 10	PRIORITY 3, 14
	LONG_INTEGER_II	Protected mode 13
	LOW LEVEL_IO 8	PUT 10
		PUT LINE 11
	Machine code insertions 7	-
		Record Clause 7
	MACHINE_EMAX 12 MACHINE_EMIN 12	Record objects 7
	MACHINE MANTISSA 12	Rendezvous 14
	MACHINE_OVERFLOWS 12	RESET II
	MACHINE_RADIX 12	Runtime Executive 1, 3, 10, 11, 13, 14
	MACHINE_ROUNDS 12	Runtime Executive limitation 15
	Main program 14	
	Main subprogram 14	SAFE_EMAX 12
	MANTISSA 12	SAFE_LARGE 12
	Maximum amount of data in the global	SAFE_SMALL 12
	data area 15	Sequential files 9
	Maximum data in the heap 15	SEQUENTIAL_IO 8, 9
	Maximum identifier length 15	SETOPT 13
	Maximum line length 15	Sharing of external files 11
	Maximum number of compilation	SHORT_INTEGER 11 SIZE 12
	units 15	JILE 14

STANDARD_INFUT 8, 10, 11 STANDARD_OUTPUT 8, 10, 11, 14 'STORAGE_SIZE 13 Storage reclamation at exit 13 Stream file 9 SUPPRESS 3 Synchronization of I/O 11 SYSTEM 3

I'SIZE 7
T'SMALL 7
T'STORAGE_SIZE 7
Tick stack size 13
Cask stacks 13
Insking issues 11
Tasks
 characteristics of 13
Text file
 buffered 10
Text files 9
TEXT_IO 8, 9, 11, 12
TIMER - FAST 14

Unchecked conversions 8
UNCHECKED_DEALLOCATION 13
Universal_real 13
Unknown media 9
USE_ERROR 9

XT without an 8087 13

APPENDIX C

TEST PARAMETERS

mentain tests in the ACVC make use of implementation-dependent values, but as the maximum length of an input line and invalid file names. A contributed makes use of such values is identified by the extension .TST in this file name. Actual values to be substituted are identified by makes that begin with a dollar sign. A value must be substituted for each of these names before the test is run. The values used for this validation are given below.

HAMP FOR MEDIUM

ed andrese served andrese andrese served departed reserved served andrese served by

Fidential or the size of the maximum input line length with varying last character.

\$BIG TD2
Identifier the size of the maximum input line length with varying last character.

\$BIG_103
Identifier the size of the
waximum input line length
with varying middle character.

Identifier the size of the maximum input line length /ith varying middle character.

\$BIG_INT_LIT
An integer literal of value 298
with enough leading zeroes so
that is is the size of the
maximum line length.

\$BIG_REAL_LIT
A real literal that can be
either of floating- or fixedpoint type, has value of 690.0,
and has enough leading zeroes to
be the size of the maximum line
length.

VALUE

NAME AND MEANING

VALUE

\$BLANKS

A sequence of blanks twenty characters fewer than the size of the maximum line length.

235 blanks

\$COURT 1051

A funiversal integer literal whose value is TEXT IO COUNTLAST.

2147483647

A string literal containing all the ASCIT characters with artitable graphics that are not a libe basic 55 Ada character

"abcdefghijklmnopqrstuvwxyz
!\$%?@[\]^'{}~"

SFIELD LAS'L

SECONDES PRESERVA DESCRIPTA DESCRIPTA DESCRIPTA DESCRIPTA DESCRIPTA DESCRIPTA DESCRIPTA DE PROPERTO DE PORTE DE

A universal integer literal whose value is TEXT IO.FIELD'LAST

255

\$FILE_NAME_WITH_BAD_CHARS
An illegal external file name
that either contains invalid
characters or is too long if no
invalid characters exist.

X}]!@ \$^&~Y

In external file name that either contains a wild card character or is too long if no wild card characters exists.

XYZ*

SGREATER_THAN_DURATION
A universal real value that lies between DURATION'BASE'LAST and DURATION'LAST if any, otherwise any value in in the range of DURATION.

2_097_151.01

\$GREATER_THAN_DURATION_BASE_LAST
The universal real value that is
greater than DURATION'BASE'LAST,
if such a value exists.

10 000 000.0

\$ILLEGAL EXTERNAL FILE NAME1
An illegal external file name.

BAD-CHARACTER*^

TEST PARAMETERS

NAME AND MEANING

VALUE

\$ILLEGAL EXTERNAL FILE NAME2
An illegal external file name
that is different from
\$ILLEGAL EXTERNAL FILE NAME1.

MUCH-TOO-LONG-NAME-FOR-A-FILE

Storeger First the universal integer literal expression whose value is INTEGERIFERST. -32768

SINTEGER LAST
The universal integer literal expression whose value is INTEGER! LAST.

32767

A universal real value that lies between DURATION'BASE'FIRST and DURATION'FIRST if any, otherwise any value in the range of DURATION.

-100 000.0

\$LESS_THAN_DURATION_BASE_FIRST
The universal real value that is
less than DURATION'BASE'FIRST, if
such a value exists.

-10 000 000.0

\$MAX_DIGITS
 The universal integer literal
 whose value is the maximum digits
 supported for floating-point types.

15

\$MAX_IN_LEN
The universal integer literal
whose value is the maximum input
line length permitted by the
implementation.

255

\$MAX_INT
 The universal integer literal
 whose value is SYSTEM.MAX INT.

2147483647

\$NAME

Tererrer (Interess () interess a factory () interess () interess () interess () interess () in the contents (

LONG_LONG_INTEGER

A name of a predefined numeric type other than FLOAT, INTEGER, SHORT_FLOAT, SHORT_INTEGER, LONG_FLOAT, or LONG_INTEGER if one exists, otherwise any undefined name.

TEST PARAMETERS

NAME AND MEANING

\$NEG_BASED_INT

A based integer literal whose highest order nonzero bit falls in the sign bit position of the representation for SYSTEM.MAX_INT.

NOT_ASCID_CHAR_TYPE
An enumerated type definition for a character type whose literals are the identifier NON_NULL and all non_ASCII characters with printable graphics.

VALUE

8#777777777776#

(NON_NULL)

WITHDRAWN TESTS

Some tests are withdrawn friche Ada Standard. The folitime of validation testing force (All-ddddd" is to an analyselfa). An unterm 8000000000000000000000000000000000000	inated string literal occurs at line ved word "IS" is misspelled at line of Election G at line 114 is ambiguace of implicit conversions. oration of subtype declarations SF2 raise NUMERIC_ERROR instead of CONST expected in the test. ct declarations at lines 126 throughbrogram bodies declared in the
Some tests are withdrawn frethe Ada Standard. The following of validation testing for the form of U-ddddd" is to an analyself4A; An unterm 8:1003C; The reserved C34018A; The call the preserved C35904A; The elab SFX4 may ERROR as a call the object of the call the call the call the preserved C3401A; The object of the call the c	on the ACVC because they do not conflowing 19 tests had been withdrawn a or the reasons indicated. A refere Ada Commentary. inated string literal occurs at line ved sord "IS" is misspelled at line of fraction G at line 114 is ambiguace of implicit conversions. oration of subtype declarations SFR raise NUMERIC_ERROR instead of CONST expected in the test. ct declarations at lines 126 througubprogram bodies declared in the
the Ada Standard. The fol- time of validation testing f the form intended. It is to an 132114A: An unterm 8 1203C: The reser C34018A: The call the prese 135904A: The elab SFX4 may ERROR as . E37401A: The obje follow s declarati C41404A: The valu the if	lowing 19 tests had been withdrawn a or the reasons indicated. A refere Ada Commentary. inated string literal occurs at line ved Gord "IS" is misspelled at line of Ametion G at line 114 is ambiguace of implicit conversions. oration of subtype declarations SFX raise NUMERIC_ERROR instead of CONST expected in the test. ct declarations at lines 126 througubprogram bodies declared in the
The reservable of the preservable of the preservabl	ved sord "IS" is misspelled at line of function G at line 114 is ambiguage of implicit conversions. oration of subtype declarations SF1 raise NUMERIC_ERROR instead of CONST expected in the test. ct declarations at lines 126 throughbrogram bodies declared in the
C34018A: The call the prese C38904A: The elab SFX4 may ERROR as C37401A: The obje follow s declaration C41404A: The valuation the if	of Ametion G at line 114 is ambiguate of implicit conversions. oration of subtype declarations SFX raise NUMERIC_ERROR instead of CONST expected in the test. ct declarations at lines 126 throughbrogram bodies declared in the
the prese The elab SFX4 may ERROR as D37401A: The obje follow s declarati C41404A: The valu the <u>if</u>	nce of implicit conversions. oration of subtype declarations SFX raise NUMERIC_ERROR instead of CONST expected in the test. ct declarations at lines 126 through upprogram bodies declared in the
SFX4 may ERROR as . 637401A: The obje follow s declarati . C41404A: The valu the <u>if</u>	raise NUMERIC_ERROR instead of CONST expected in the test. ct declarations at lines 126 through ubprogram bodies declared in the
follow s declarati C41404A: The valu the <u>if</u>	ubprogram bodies declared in the
the <u>if</u>	
	es of 'LAST and 'LENGTH are incorrestatements from line 74 to the end of
value of	<pre>l and ARRPRIBL 2 are initialized v the wrong typePRIBOOL TYPE inste L_TYPEat line 41.</pre>
values o	mption that evaluation of default cours when an exception is raised is incorrect according to AI-00397.
terminate	declarations at lines 41 and 50 d incorrectly with colons, and end g from line 42.
	ect declaration in line 18 foll m body of the same declarative part.
	D-1

WITHDRAWN TESTS

. B74101B: The <u>begin</u> at line 9 causes a declarative part to be treated as a sequence of statements.

. C87B50A: The call of "/=" at line 31 requires a use clause for package A.

C92005A: The "/=" for type PACK.BIG_INT at line 40 is not visible without a use clause for the package PACK.

The assumption that allocated task TT1 will run prior to the main program, and thus assign SPYNUMB the value checked for by the main program, is erroneous.

No valid elaboration order exists for these tests.

The body of BC3204C0 is missing.

Secretary described property and property of the secretary

L M D DATE FILMED 8-8 DTIC